# A System to Detect Forged-Origin BGP Hijacks

### **Thomas Holterbach** University of Strasbourg

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### Joint work with: Alberto Dainotti Thomas Alfroy Amreesh D. Phokeer **Cristel Pelsser**



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Autonomous System (AS)





path: 3 4 7







### ASes that divert the traffic to 7.0.0/8 to the attacker





# Fortunately, there are defenses against BGP hijacking

Protocol extensions

RPKI + ROV BGPSec, ASPA

Configuration guidelines

**Route filters** 

Monitoring platforms

ARTEMIS BGPAlerter

# Despite the efforts, BGP is *still* vulnerable to forged-origin hijacks

The attacker prepends the legitimate AS number to the AS path



# Despite the efforts, BGP is *still* vulnerable to forged-origin hijacks

### Less but still a significant fraction of the traffic is diverted to the attacker





# Existing defenses poorly neutralise forged-origin hijacks

### **Protocol** extensions

### **RPKI + ROV** BGPSec, ASPA

### Configuration guidelines

### **Route filters**

### Monitoring platforms

ARTEMIS **BGPAlerter** 



**RPKI+ROV** can't detect forged-origin hijacks **BGPSec and ASPA will take years** to be widely deployed

Often missing and inaccurate as they are constructed based on the IRR

Narrowly focused as they detect hijacks that only pertain to the AS deploying it



# Forged-origin hijacks are actively used by attackers



February 3, 2022

### KlaySwap crypto users lose funds after BGP hijack

Hackers have stolen roughly \$1.9 million from South Korean cryptocurrency platform KLAYswap after they pulled off a rare and clever BGP hijack against the server infrastructure of one of the platform's providers.

The BGP hijack—which is the equivalent of hackers hijacking internet routes to bring users on malicious sites instead of legitimate ones—hit KakaoTalk, an instant messaging platform popular in South Korea.

The attack took place earlier this month, on February 3, lasted only for two hours, and KLAYswap has confirmed the incident last week and is currently issuing compensation for affected users.

### Both attacks are the result of a forged-origin hijack

### August 17, 2022





# **DFOH:** A System to Detect Forged-Origin BGP Hijacks on the Whole Internet





### **DFOH's main challenge**

### **DFOH's inference pipeline**

### **DFOH's inferences are accurate**

**DFOH** is up and running



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**DFOH's inference pipeline** 

**DFOH's inferences are accurate** 

**DFOH** is up and running

### is to detect fake AS links

### **DFOH** aims to detect the fake AS links induced by forged-origin hijacks BGP vantage point $\mathbf{\Theta}\mathbf{\Theta}$ path: 3 5 7 00 path: 4 7 2 Traffic to . . . . . . . . . . . . . . . . 7.0.0/8 path: 7 path: 5 7 3 4 6 **Attacker:** announces path: 7 7.0.0/8 hijacks 7.0.0/8 prepends 7 15



### **DFOH** aims to detect the fake AS links induced by forged-origin hijacks BGP vantage point $(\mathbf{0})$ path: 3 5 7 •• path: 4 7 2 Traffic to Upon the attack: . . . . . . . . . . . . . . 7.0.0/8 AS5 (attacker) and AS7 (victim) appear directly connected \* path: 7 path: 5 7 3 4 6 \* **\* Attacker:** announces path: 7 7.0.0/8 hijacks 7.0.0/8

16

fake link

prepends 7



# An attacker cannot escape from creating a fake AS link without hampering the effectiveness of its attack

There is no new AS link if the attacker prepends 67

But none of the ASes divert traffic to the attacker as the AS path is longer

> **Attacker:** hijacks 7.0.0/8 prepends 67





# **Problem:** There are many new AS links every day but no simple property that tells whether they are real or fake



We find 166 new AS links every day (median) and the vast majority are likely legitimate

Using the BGP data from 200 RIS and RouteViews peers and collected during ten months in 2022



# <u>Problem:</u> There are many new AS links every day but no simple property that tells whether they are real or fake







### **DFOH's main challenge**

### **DFOH's inference pipeline** relies on domain-specific knowledge and a tailored link prediction framework

### **DFOH's inferences are accurate**

**DFOH** is up and running

### is to detect fake AS links









**Feature vectors** 

# **DFOH** uses a total of **11** topological features that can be divided into four categories

Node centrality Neighborhood richness





### Topological patterns

### Closeness

focus



shortest distance





**Feature vectors** 

# **DFOH** looks at public peering information and identifies when two ASes are unlikely to peer

**DFOH** looks for three types of information in PeeringDB:

- 1. Country
- 2. Public peering exchange points
- 3. Private peering facilities

# **DFOH** compares the peering information of the neighbors of the hypothetical victim and attacker

### Reason #1: Protect against adversarial inputs



Reason #2: Mitigate missing peering information

![](_page_27_Figure_1.jpeg)

**Feature vectors** 

# **DFOH** looks at the AS paths that include the new link and identifies suspicious sequence of ASes

### AS node degree

![](_page_28_Figure_2.jpeg)

# **DFOH** looks at the AS paths that include the new link and identifies suspicious sequence of ASes

![](_page_29_Figure_1.jpeg)

### AS node degree

# **DFOH** looks at the AS paths that include the new link and identifies suspicious sequence of ASes

![](_page_30_Figure_1.jpeg)

![](_page_30_Figure_2.jpeg)

stub hijacks core AS

![](_page_30_Figure_4.jpeg)

![](_page_31_Figure_1.jpeg)

![](_page_32_Figure_1.jpeg)

![](_page_33_Figure_1.jpeg)

![](_page_33_Picture_3.jpeg)

![](_page_33_Figure_4.jpeg)

# There are several link prediction frameworks SEAL (NIPS'18) is one example

![](_page_34_Picture_1.jpeg)

# There are several link prediction frameworks but they do not translate well for detecting fake AS links

### Few tier1 ASes

### Many stub ASes

![](_page_35_Picture_3.jpeg)

Clusters of ASes based on their degree and cone size

### Stub

Transj.

- Transit/IXP/CDN 1
- Transit/IXP/CDN 2 -
- Transit/IXP/CDN 3 -
- Transit/IXP/CDN 4 -
- Highly connected
- Large customer cone -
  - Tier1

![](_page_36_Figure_10.jpeg)

![](_page_36_Figure_11.jpeg)

![](_page_36_Picture_13.jpeg)

![](_page_37_Figure_1.jpeg)

Stub 98%

Transj.

Transit/IXP/CDN 1

Transit/IXP/CDN 2

Transit/IXP/CDN 3

Transit/IXP/CDN 4

Highly connected

Large customer cone

Tier1

![](_page_37_Figure_10.jpeg)

Proportion of sampled **nonexistent** AS links (random sampling)

2%

![](_page_37_Picture_14.jpeg)

Clusters of ASes based on their degree and cone size

**DFOH would perform** well on scenarios involving two stubs

Transit/IXP/CDN 1

Transit/IXP/CDN 2

Transit/IXP/CDN 3

Transit/IXP/CDN 4

Highly connected

Large customer cone

Tier1

Transj+

Stub

![](_page_38_Figure_10.jpeg)

![](_page_38_Picture_12.jpeg)

![](_page_39_Figure_1.jpeg)

![](_page_39_Picture_3.jpeg)

![](_page_40_Figure_1.jpeg)

![](_page_40_Picture_3.jpeg)

![](_page_40_Figure_4.jpeg)

![](_page_41_Picture_0.jpeg)

### **DFOH's main challenge**

# **DFOH's inference pipeline**

### **DFOH**'s inferences are accurate

### **DFOH** is up and running

### is to detect fake AS links

relies on domain-specific knowledge and a tailored link prediction framework

in every attack scenario

# We evaluate **DFOH** on artificially created forged-origin hijacks as there is no ground truth at scale

**Methodology:** 

We take existing AS paths and prepend a new origin to create a new link

We take 9k cases where the new link exists (legitimate or "negative" cases) and 9k cases where the new link does not exist (suspicious or "positive" cases)

![](_page_42_Picture_5.jpeg)

# We evaluate **DFOH** on artificially created forged-origin hijacks as there is no ground truth at scale

**Methodology:** 

We take existing AS paths and prepend a new origin to create a new link

We take 9k cases where the new link exists (legitimate or "negative" cases) and 9k cases where the new link does not exist (suspicious or "positive" cases)

![](_page_43_Picture_4.jpeg)

We focus on the True Positive Rate (TPR) and the False Positive Rate (FPR)

![](_page_43_Picture_8.jpeg)

# **True Positive** Rate

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Attac

Transitivity it is 'to CON 3

Stub

Transit/IXP/CDN 1 -

Transit/IXP/CDN 2 -

Transit/IXP/CDN 3 -

Transit/IXP/CDN 4 -

Highly connected

Large customer cone -

Tier1

### Victim

![](_page_44_Picture_13.jpeg)

![](_page_44_Picture_14.jpeg)

# **True Positive** Rate

![](_page_45_Picture_3.jpeg)

Attacker

Stub -	0.97	0.86	0.91	0.96
Transit/IXP/CDN 1	0.86	0.73	0.90	0.97
Transit/IXP/CDN 2	0.91	0.90	0.85	0.95
Transit/IXP/CDN 3	0.96	0.97	0.95	0.99
Transit/IXP/CDN 4	0.94	0.82	0.99	1.00
Highly connected	0.95	0.96	0.99	0.98
Large customer cone	0.95	0.83	0.90	0.99
Tier1	0.84	0.73	0.83	0.91

### Victim

46

![](_page_45_Picture_7.jpeg)

# **True Positive** Rate

![](_page_46_Picture_3.jpeg)

Attacker

La

0.97	0.86	0.91	0.96
0.86	0.73	0.90	0.97
0.91	0.90	0.85	0.95
0.96	0.97	0.95	0.99
0.94	0.82	0.99	1.00
0.95	0.96	0.99	0.98
0.95	0.83	0.90	0.99
0.84	0.73	0.83	0.91
	0.97 0.86 0.91 0.96 0.94 0.95 0.95	0.970.860.860.730.910.900.960.970.950.820.950.830.840.73	<table-container>  0.977 0.866 0.911   0.866 0.733 0.900   0.911 0.900 0.855   0.926 0.927 0.995   0.955 0.832 0.900   0.841 0.833 0.900</table-container>

### Victim

![](_page_46_Picture_7.jpeg)

### The minimum TPR is 0.73

# **False Positive** Rate

Attacker

![](_page_47_Picture_3.jpeg)

Stub

Transit/IXP/CDN 1

Transit/IXP/CDN 2 -

Transit/IXP/CDN 3

Transit/IXP/CDN 4

Highly connected

Large customer cone

Tier1

## Victim

![](_page_47_Picture_13.jpeg)

![](_page_48_Picture_1.jpeg)

# **False Pos** Rate

![](_page_48_Picture_3.jpeg)

Attacker

ositive te	Transit/IXp Stu	Trans situry CON	Trans it it it it is			ASton	, , , , , , , , , , , , , , , , , , ,	$T_{i_{\Theta_{r}}}$	7
	Stub	- 0.04	0.03	0.02	0.01	0.00	0.01	0.02	0.03
Transit/IX	P/CDN 1	- 0.03	0.03	0.01	0.01	0.02	0.00	0.02	0.06
Transit/IX	P/CDN 2	- 0.02	0.01	0.02	0.01	0.03	0.01	0.03	0.07
Transit/IX	P/CDN 3	- 0.01	0.01	0.01	0.00	0.05	0.01	0.03	0.00
Transit/IX	P/CDN 4	- 0.00	0.02	0.03	0.05	0.04	0.01	0.00	0.06
Highly co	onnected	- 0.01	0.00	0.01	0.01	0.01	0.00	0.00	0.15
Large custor	ner cone	- 0.02	0.02	0.03	0.03	0.00	0.00	0.03	0.07
	Tier1	- 0.03	0.06	0.07	0.00	0.06	0.15	0.07	0.02
						49			

### Victim

![](_page_48_Picture_7.jpeg)

![](_page_49_Picture_1.jpeg)

# **False Positive** Rate

![](_page_49_Picture_3.jpeg)

Attacker

		-	-	=
Stub	- 0.04	0.03	0.02	0.01
Transit/IXP/CDN 1	- 0.03	0.03	0.01	0.01
Transit/IXP/CDN 2	- 0.02	0.01	0.02	0.01
Transit/IXP/CDN 3	- 0.01	0.01	0.01	0.00
Transit/IXP/CDN 4	- 0.00	0.02	0.03	0.05
Highly connected	- 0.01	0.00	0.01	0.01
Large customer cone	- 0.02	0.02	0.03	0.03
Tier1	- 0.03	0.06	0.07	0.00

### Victim

![](_page_49_Picture_7.jpeg)

### The maximum FPR is 0.15

![](_page_50_Picture_0.jpeg)

### **DFOH's main challenge**

### **DFOH's inference pipeline discriminates** fake AS links from the real ones

### **DFOH's inferences are accurate** in

### **DFOH** is up and running

and useful for operators

### is to detect fake AS links

in every attack scenario

### **DFOH** runs at https://dfoh.uclouvain.be

### DFOH

### A System to Detect Forged-Origin BGP Hijacks

DFOH is a system that aims to detect forged-origin hijacks on the whole Internet. Forged-origin hijacks are a type of BGP hijack where the attacker manipulates the AS path of BGP messages to make them appear as legitimate routing updates.

**DFOH** is useful given that the BGP extensions proposed to cryptographically verify the validity of the AS paths (such as BGPSec or ASPA) are hard to widely deploy. With DFOH, operators can quickly and with high confidence know when their IP prefixes are being hijacked.

### Read our NSDI'24 paper

### A System to Detect Forged-Origin BGP Hijacks

Thomas Holterbach," Thomas Alfroy," Amreesh Phokeer," Alberto Dainotti, Cristel Pelsse \*University of Strasbourg, <sup>1</sup>Internet Society, <sup>1</sup>Georgia Tech, <sup>3</sup>UCLouvain

oit the lack of strong BGP securi ijacks, a type of BGP hijac s the AS path to make it i route is legitimate or has been manymiated. We demonstrate that current state-for at approaches to direct: BGP anoenalies are insufficient to deal with forged-origin hijacks. We identify the key properties that make the inference of forged AS paths challenging, and design DFOP to be robust against real-world factors (e.g., data biases). Our inference pipeline includes two key ingredients: (i) a set of strategically selected features, and (ii) a training scheme adapted to topological biases. DFOH detects 90.9% of the forged-origin hijacks within only <5min. In addition, in only reports <7.7.5 suspicious cease severy day for the whole Internet, a small number that allows operators to investigate the reported cases and take countermeasures.

### 1 Introduction

On 3 February 2022, the cryptocarrency platform KLAYswap was targeted by hackers who stole \$1.9 million worth of digi-tal assets [59]. More recently, on 17 August 2022, an attack This assets [59]. More recently, on 17 August 2022, in matrix, to Bridge-accentred S2 victims, where the stacks lost 2325,000 [4]. Both attacks were the result of a forged-origin BGP diack, a type of routing higher. Where the attacks ers announce foregot AS paths towards a victim prefix by prepending the victim's origin AS number in order to make them appear legitimate. Clearly, BGP higher attacks are and a forged-origin highers. In a surgive anymore: They repeatedly make the headlines [1,2] and are known us attack vertors to sited cryptocurrency [8], obtain bogus certificates [15], or deanonymize Tor users [62].

to detect forged-origin hijacks, since the forged origin in the AS path is actually valid, while IRR-based filters are known As path is actually valid, while IRR-based filters are known to be inaccurate, incomplete [25], and far two often missing given the increasing number of observed BGP hijacks [7]. Today, network operators do not have many options left other than widing for the deployment of new security extensions to BGP to consistently prevent forged-crigin hijacks [44]. Such deployment—II is will happen at all—might take more than a decade, as in the case of RFK1-ROV [21]. In this paper, we present DFCH, the first *locally-deployable* system that wide'r, *owdely*, and *accurate* b heters Forced

system that widely, quickly, and accurately Detects Forged-Origin Hijacks on the Internet. With a single deployment of DFOH on a commodity server, any attacker performing a forged-origin hijack is likely to be quickly detected, the hijack publicly reported, and the victim immediately notified. Being publicity reported, and the victim immediately notified. Being aware of the attack, the victim can apply countermeasure and the community can take actions to prevent similar attack from happening again. Additionally, DFOM can detect pas attacks, allowing the community to measure the frequency or

### Watch our APRICOT'24 presentation

![](_page_51_Picture_16.jpeg)

### **DFOH** provides past and real-time forged-origin BGP hijacks detection

![](_page_51_Picture_19.jpeg)

## **DFOH** is useful and practical for network operators

# <u>Useful:</u> DFOH detects the two known forged-origin BGP hijacks (the klayswap and cbridge attacks)

Practical: DFOH only reports zero or o (worse case is 15 cases)

### Practical: DFOH only reports zero or one case every month for 99.8% of the ASes

# **DFOH:** A System to Detect Forged-Origin Hijacks

![](_page_53_Picture_1.jpeg)

# *DFOH* runs in a commodity server

![](_page_53_Picture_3.jpeg)

**DFOH** detects hijacks on the whole Internet

CDN Tier1 Stub

**DFOH** is accurate in every attack scenario

# **DFOH:** A System to Detect Forged-Origin BGP Hijacks

![](_page_54_Picture_1.jpeg)

# *DFOH* runs in a commodity server

![](_page_54_Picture_3.jpeg)

**DFOH** detects hijacks on the whole Internet

CDN Tier1 Stub

**DFOH** is accurate in every attack scenario

![](_page_54_Picture_7.jpeg)

**DFOH** detects past hijacks

![](_page_54_Picture_9.jpeg)

**DFOH** provides near-real-time detection

![](_page_54_Picture_11.jpeg)

**DFOH** is robust against adversarial inputs

# **DFOH:** A System to Detect Forged-Origin Hijacks https://dfoh.uclouvain.be

![](_page_55_Picture_1.jpeg)

# **DFOH** runs in a commodity server

![](_page_55_Picture_3.jpeg)

**DFOH** detects hijacks on the whole Internet

![](_page_55_Figure_5.jpeg)

**DFOH** is accurate in every attack scenario

![](_page_55_Picture_7.jpeg)

![](_page_55_Picture_8.jpeg)

**DFOH** detects past hijacks

![](_page_55_Picture_10.jpeg)

**DFOH** provides near-real-time detection

![](_page_55_Picture_12.jpeg)

**DFOH** is robust against adversarial inputs