### A System to Detect Forged-Origin Hijacks

#### **Thomas Holterbach University of Strasbourg**

#### **Routing Security Summit** 2023

#### Joint work with: Thomas Alfroy Alberto Dainotti Amreesh D. Phokeer **Cristel Pelsser**















#### Fortunately, there are defences against BGP hijacking

### Protocol extensions

#### RPKI + ROV BGPSec, ASPA

### Configuration guidelines

#### **Route filters**

#### Monitoring platforms

ARTEMIS BGPAlerter

### Despite the efforts, BGP is *still* vulnerable to forged-origin hijacks





#### Existing defenses poorly neutralise forged-origin hijacks

#### **Protocol** extensions

#### **RPKI + ROV BGPSec, ASPA**

#### Configuration guidelines

#### **Route filters**

#### Monitoring platforms





**RPKI+ROV** can't detect forged-origin hijacks ASPA will take years to be deployed







#### Forged-origin hijacks are actively used by attackers



February 3, 2022

#### KlaySwap crypto users lose funds after BGP hijack

Hackers have stolen roughly \$1.9 million from South Korean cryptocurrency platform KLAYswap after they pulled off a rare and clever BGP hijack against the server infrastructure of one of the platform's providers.

The BGP hijack—which is the equivalent of hackers hijacking internet routes to bring users on malicious sites instead of legitimate ones—hit KakaoTalk, an instant messaging platform popular in South Korea.

The attack took place earlier this month, on February 3, lasted only for two hours, and KLAYswap has confirmed the incident last week and is currently issuing compensation for affected users.

#### Both attacks are the result of a forged-origin hijack

#### August 17, 2022





### **DFOH:** A System to Detect Forged-Origin Hijacks on the Whole Internet

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#### Outline

#### 1. *DFOH*'s main challenge is to detect fake AS links

#### 2. **DFOH's key ingredients are** carefully selected features and a balanced sampling

3. **DFOH** is accurate and practical for users

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### **DFOH** aims to detect the fake AS links induced by forged-origin hijacks





### **DFOH** aims to detect the fake AS links induced by forged-origin hijacks

BGP vantage point





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BGP vantage point





### An attacker cannot escape from creating a new AS link without hampering the effectiveness of its attack





### **<u>Problem</u>:** There are many new AS links every day but no simple property that tells whether they are real or fake



We find 166 new AS links every day (median)

Using the BGP data from 200 RIS and RouteViews peers and collected during ten months in 2022



# <u>Problem:</u> There are many new AS links every day but no simple property that tells whether they are real or fake







#### 1. **DFOH**'s main challenge is to detect fake AS links

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**Feature vectors** 

### **DFOH** uses a total of **11** topological features that can be divided into four categories

Node centrality Neighborhood richness





#### Topological patterns

#### Closeness









**Feature vectors** 

#### **DFOH** leverages correlations in the public peering information

**DFOH** looks for three types of information in PeeringDB:

1. Country

2. Public peering exchange points

3. Private peering facilities



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EQUINIX





**Feature vectors** 

# **DFOH** detects fake AS paths as they often violate patterns induced by business relationships



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## towards stub-to-stub links as they are overrepresented

**Problem:** randomly sampling nonexistent links makes DFOH skewed



Clusters of ASes based on their degree and cone size

#### Stub

Transj.

- Transit/IXP/CDN 1
- Transit/IXP/CDN 2 -
- Transit/IXP/CDN 3 -
- Transit/IXP/CDN 4 -
- Highly connected
- Large customer cone -
  - Tier1



Proportion of sampled **nonexistent** AS links (random sampling)









DFOH would perform < well on scenarios involving two stubs

Transit/IXP/CDN 1 - 0.02

Stub

Transit/IXP/CDN 2 - 0.00

Transit/IXP/CDN 3 - 0.00

Transit/IXP/CDN 4

Highly connected - 0.00

Tier1



Proportion of sampled **nonexistent** AS links (random sampling)





0.02	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00
0.00	0.00	0.00	0.00	0.00	0.00	0.00

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### 1. *DFOH*'s main challenge is to detect fake AS links

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3. *DFOH* is accurate and practical for users

#### We evaluate **DFOH** on artificially created forged-origin hijacks and measure its accuracy upon every attack scenario

Methodology:

Step #1: We take existing AS paths and prepend a new origin to create a new link

Step #2: We consider 9k cases where the new link exists (*legitimate cases*) and 9k cases where the new link does not existent (*malicious cases*)

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We focus on the True Positive Rate (TPR) and the False Positive Rate (FPR)

Victim



Stub -	0.97	0.86	0.91	0.96	0.94	0.95	С
Transit/IXP/CDN 1	0.86	0.73	0.90	0.97	0.82	0.96	С
Transit/IXP/CDN 2	0.91	0.90	0.85	0.95	0.99	0.99	C
Transit/IXP/CDN 3	0.96	0.97	0.95	0.99	1.00	0.98	С
Transit/IXP/CDN 4	0.94	0.82	0.99	1.00	0.90	1.00	С
Highly connected	0.95	0.96	0.99	0.98	1.00	1.00	1
Large customer cone	0.95	0.83	0.90	0.99	0.85	1.00	C
Tier1	0.84	0.73	0.83	0.91	0.83	0.96	С

TPR

Attacker



Victim



0.97 0.86 0.91 0.96 0.94 0.95 0.95 0.84 Stub Transit/IXP/CDN 1 - 0.86 0.73 0.90 0.97 0.82 0.96 0.83 0.73 0.91 0.90 0.85 0.95 0.99 0.99 Transit/IXP/CDN 2 0.96 0.97 0.95 0.99 1.00 0.98 0.99 0.91 Transit/IXP/CDN 3 0.94 0.82 0.99 1.00 0.90 1.00 0.85 0.83 Transit/IXP/CDN 4 0.95 0.96 0.99 0.98 1.00 1.00 1.00 0.96 Highly connected 0.95 0.83 0.90 0.99 0.85 1.00 0.97 0.89 Large customer cone 0.84 0.73 0.83 0.91 0.83 0.96 0.89 0.78 Tier1

TPR

Attacker



#### The minimum TPR is 0.73





FPR
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Attacker

Stub	- 0.04	0.03	0.02	0.01	0.00	0.01	(
Transit/IXP/CDN 1	- 0.03	0.03	0.01	0.01	0.02	0.00	C
Transit/IXP/CDN 2	- 0.02	0.01	0.02	0.01	0.03	0.01	(
Transit/IXP/CDN 3	- 0.01	0.01	0.01	0.00	0.05	0.01	(
Transit/IXP/CDN 4	- 0.00	0.02	0.03	0.05	0.04	0.01	(
Highly connected	- 0.01	0.00	0.01	0.01	0.01	0.00	(
arge customer cone	- 0.02	0.02	0.03	0.03	0.00	0.00	(
Tier1	- 0.03	0.06	0.07	0.00	0.06	0.15	(







Transit/IXP/CDN 1 - 0.03 0.03 0.01 0.01 0.02 0.00 0.02 0.06 Transit/IXP/CDN 2 - 0.02 0.01 0.02 0.01 0.03 0.01 0.03 0.07 Transit/IXP/CDN 3 - 0.01 0.01 0.01 0.00 0.05 0.01 0.03 0.00 Transit/IXP/CDN 4 - 0.00 0.02 0.03 0.05 0.04 0.01 0.00 0.06 

**FPR** 

**ttacke** 

## **DFOH** makes the detection of forged-origin hijacks practical for operators



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### **DFOH:** A System to Detect Forged-Origin Hijacks



## *DFOH* runs in a commodity server



**DFOH** detects hijacks on the whole Internet

CDN Tier1 Stub

**DFOH** is accurate in every attack scenario

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### **DFOH** detects past hijacks



**DFOH** provides near-real-time detection



**DFOH** is robust against adversarial inputs



# **DFOH:** A System to Detect Forged-Origin Hijacks dfoh.info.ucl.ac.be



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## An attacker cannot escape from creating a new AS link without hampering the effectiveness of its attack



### An attacker cannot escape from creating a new AS link without hampering the effectiveness of its attack



prepended by the attacker

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% of polluted ASes

prepended by the attacker



6





Pattern #1: Hierarchical structure with a few Tier1s and many stubs



~70k stub ASes



Pattern #1: Hierarchical structure with a few Tier1s and many stubs

Pattern #2: CDNs and HyperGiants are highly connected



Pattern #1: Hierarchical structure with a few Tier1s and many stubs

Pattern #2: CDNs and HyperGiants are highly connected

Pattern #3: Remote peerings and IP tunnels flatten the graph



![](_page_57_Picture_5.jpeg)

#### <u>Step #1:</u> Finding new links

#### **DFOH** takes all updates and one RIB per month from 200 BGP vantage points selected using MVP\*

**DFOH** builds the AS topology at day d using AS paths in BGP routes collected during the 300 days prior d

**DFOH** infers that an AS link observed at day d is new if the link is not in the AS topology constructed at day d

\*mvp.info.ucl.ac.be

![](_page_58_Picture_7.jpeg)

## **DFOH** computes the change induced by the new AS link on topological features

3

Example with the shortest distance feature

Before the new link: shortest distance between 6 and 9 is 5

After the new link: shortest distance between 6 and 9 is 1

**Difference** is 4

6

![](_page_59_Figure_6.jpeg)

New AS link

![](_page_59_Picture_8.jpeg)

# **DFOH** verifies that a new AS link is observed in both directions as it is a strong indicator of legitimacy

- **DFOH** verifies link bidirectionality using:
- BGP data (AS path) from many VPs
- IRR data (import/export policies)

![](_page_60_Figure_4.jpeg)

![](_page_60_Figure_5.jpeg)

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![](_page_61_Figure_4.jpeg)

![](_page_61_Figure_5.jpeg)

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- **DFOH** verifies link bidirectionality using:
- BGP data (AS path) from many VPs
- IRR data (import/export policies)

AS links observed in both directions are legitimate

![](_page_62_Picture_5.jpeg)

3

![](_page_62_Figure_6.jpeg)

#### The bidirectionality feature is safe as an attacker cannot intentionally fake both directions of an AS link

Faking both directions in the same AS path would create a loop

Faking both directions in the IRR is not possible as the attacker only controls its IRR data

Merging these two datasets is safe as an attacker can only fake the same direction in BGP and the IRR

# **DFOH** considers the neighbouring nodes to avoid adversarial inputs as the information on Peeringdb is not verified

![](_page_64_Figure_1.jpeg)

New AS link

# **DFOH** learns the pattern of legitimate and malicious AS paths using a supervised training model

**DFOH** samples X legitimate AS paths and artificially creates the same number of maliciously-induced AS paths

**DFOH** computes the degree and customer cone size of every AS in the sampled AS paths

**DFOH** trains a random forest that it uses to compute a probability that a given AS path is fake or real

### **DFOH** builds a sample of nonexistent links that is similarly balanced as the set of existing links

#### **Existing links distribution** between different AS categories

AS category:  $\rightarrow$ 

	1	I	1	1	1	- 1	1	1	Г	0.12
_	0.13	0.16	0.11	0.14	0.02	0.10	0.04	0.09		
_	0.16	0.05	0.06	0.04	0.01	0.02	0.01	0.01	-	0.10
_	0.11	0.06	0.01	0.01	0.00	0.00	0.00	0.00	-	0.08
_	0.14	0.04	0.01	0.00	0.00	0.00	0.00	0.00		0.06
_	0.02	0.01	0.00	0.00	0.00	0.00	0.00	0.00		0.00
_	0.10	0.02	0.00	0.00	0.00	0.00	0.00	0.00	F	0.04
_	0.04	0.01	0.00	0.00	0.00	0.00	0.00	0.00	ŀ	0.02
_	0.09	0.01	0.00	0.00	0.00	0.00	0.00	0.00		0.00
										0.00

#### **Sampled nonexistent links distribution** when using **DFOH**'s balanced sampling

![](_page_66_Figure_5.jpeg)

![](_page_66_Figure_6.jpeg)

# **DFOH** uses a random forest classifier to classify an AS link as fake or legitimate

1. DFOH samples 30k existing and nonexistent AS links

2. DFOH estimates the best parameters using a cross-validated grid search on 25% of the sampled AS links

3. *DFOH* trains the classifier with the remaining 75% of the AS links

![](_page_67_Picture_4.jpeg)

**DFOH** repeats this process every day to ensure that its inferences remain accurate over time